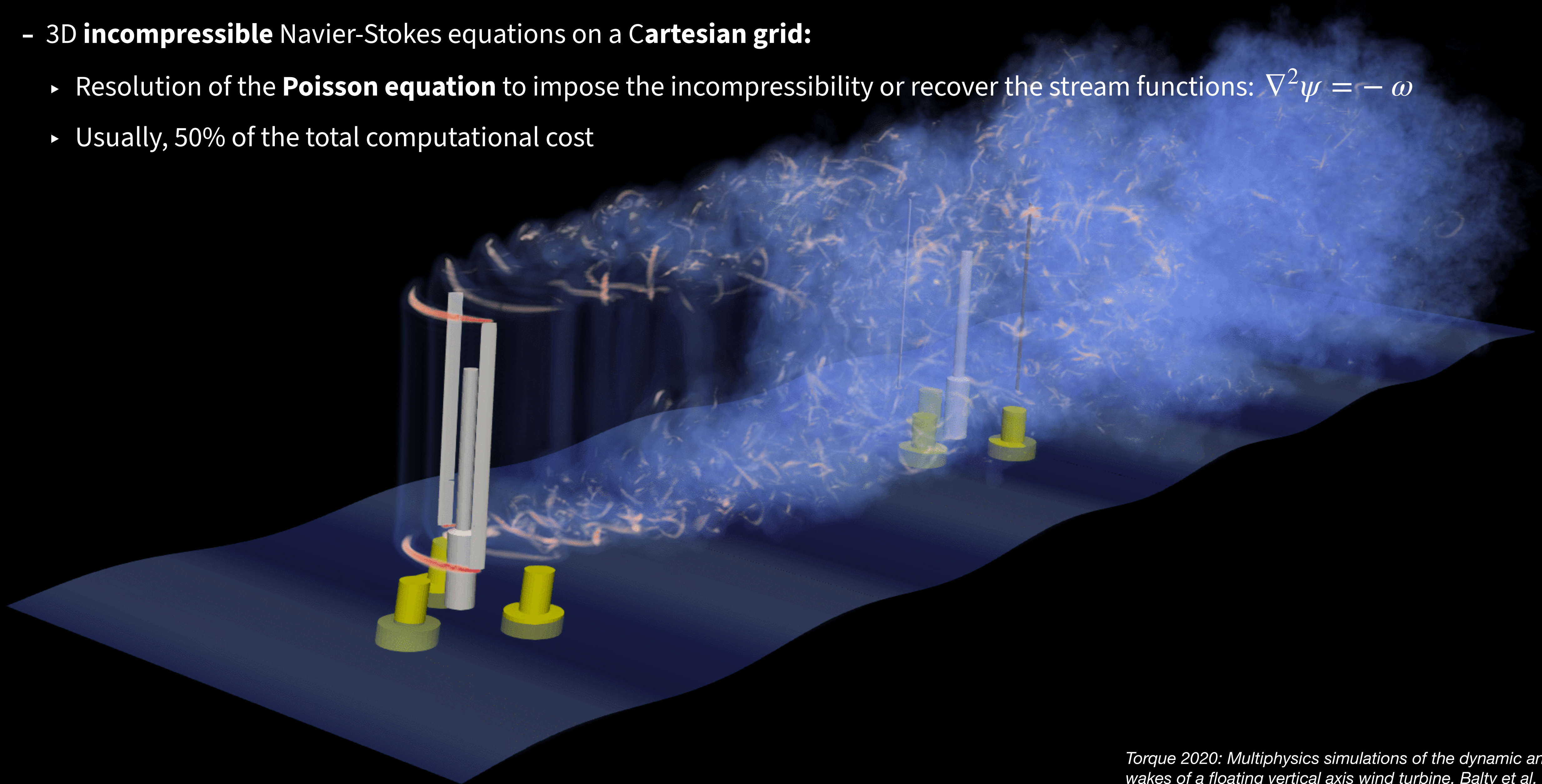


FLUPS: a versatile and performant FFT-based library of unbounded Poisson solvers

P. Balty, D.-G. Caprace, P. Chatelain, T. Gillis

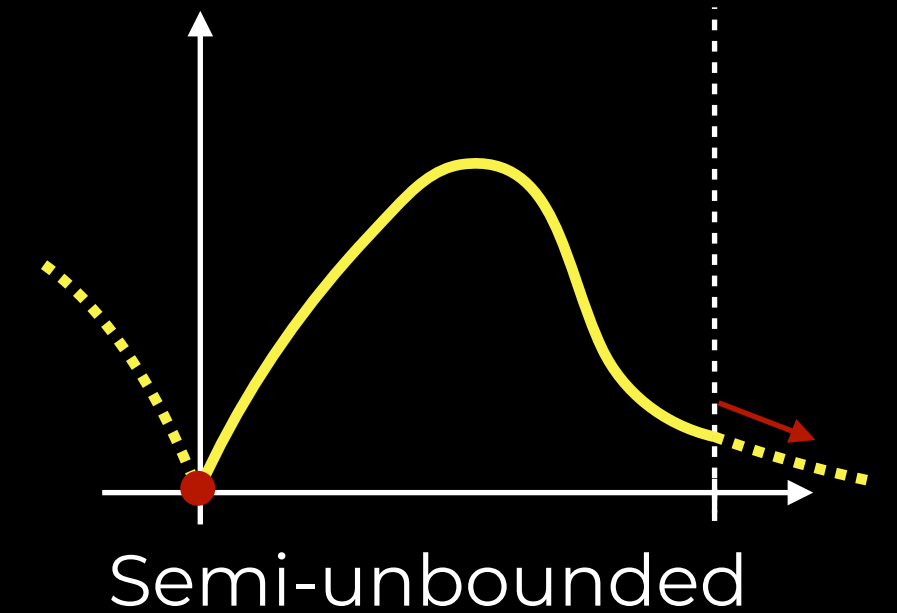
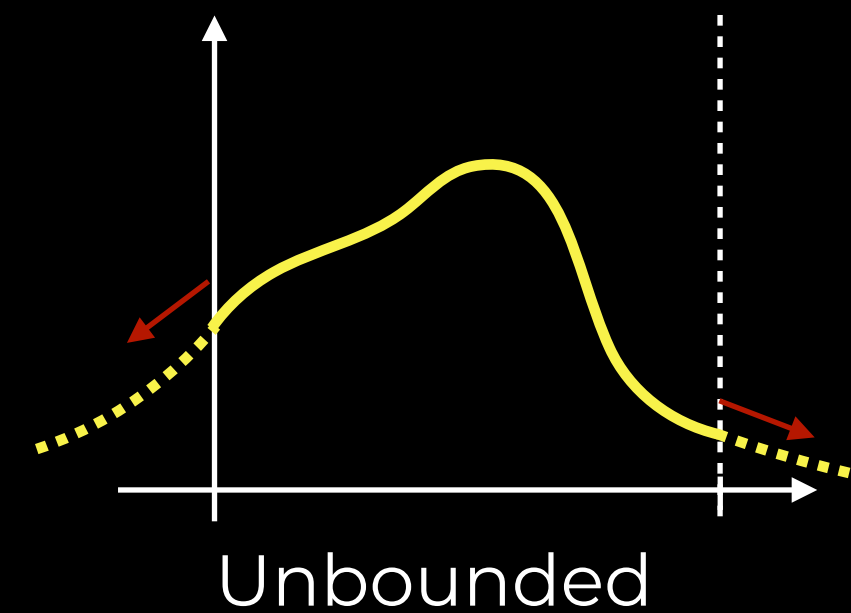
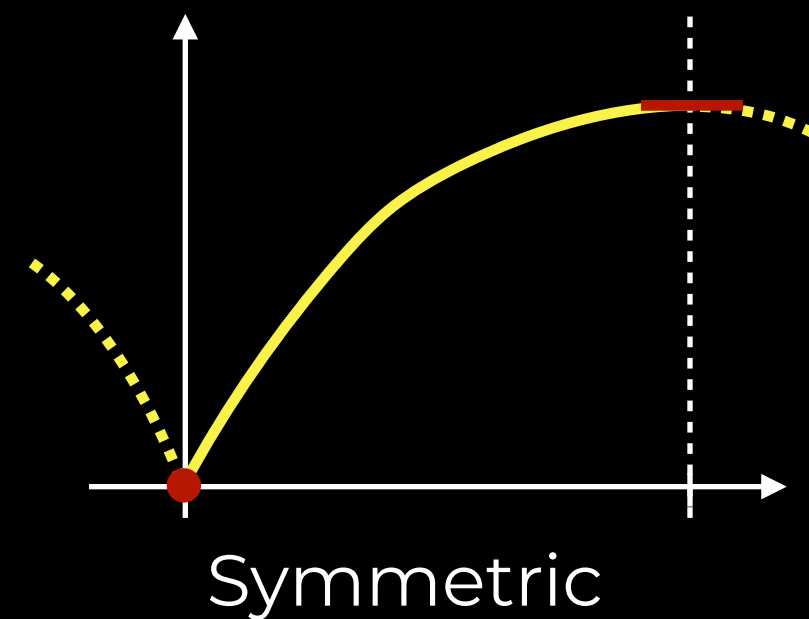
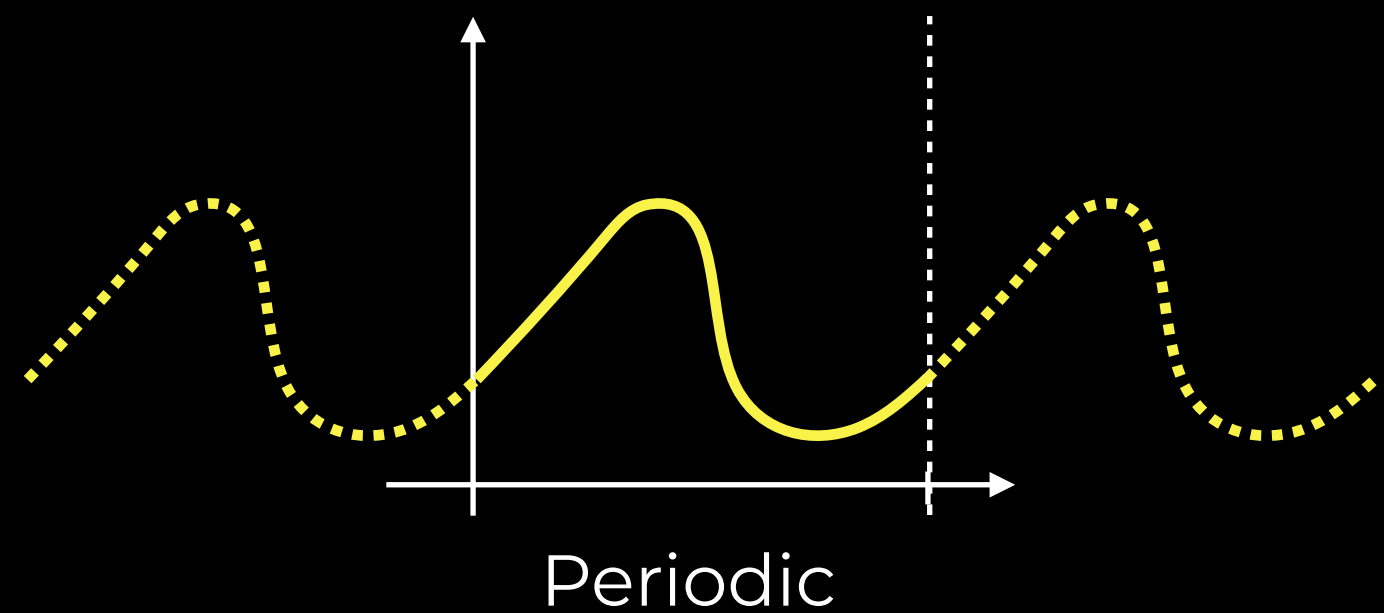
Context

- 3D **incompressible** Navier-Stokes equations on a **Cartesian grid**:
 - Resolution of the **Poisson equation** to impose the incompressibility or recover the stream functions: $\nabla^2 \psi = -\omega$
 - Usually, 50% of the total computational cost



Context

- 3D **incompressible** Navier-Stokes equations on a **Cartesian grid**:
 - Resolution of the **Poisson equation** to impose the incompressibility or recover the stream functions: $\nabla^2 \psi = -\omega$
 - Usually, 50% of the total computational cost
- Specific need:
 - Performance on massively parallel systems
 - Flexibility in the data layout: cell-centered or node-centered
 - Various boundary conditions (BC):



- FFT-based method is compatible with most of the BCs and are the fastest on uniform rectangular grid ^[1].

Plan for today

- I. Context
- II. Methodology
- III. Implementation for massively distributed systems
- IV. Results
- V. Conclusion

An unbounded FFT-based Poisson Solver...

- FFT-based method:

$$\nabla^2 \Psi = -f \longrightarrow \psi = G \star f \longrightarrow \hat{\psi} = \hat{G} \hat{f}$$

- 3D FFT obtained as a succession of 1D FFTs: $\psi_{x,y,z} \rightarrow \tilde{\psi}_{x,ky,z} \rightarrow \tilde{\psi}_{x,ky,kz} \rightarrow \hat{\psi}_{kx,ky,kz}$
- Analytical expression of the Green's functions (provided some regularisation of a given order)
- Different transforms for each boundary condition and data layout

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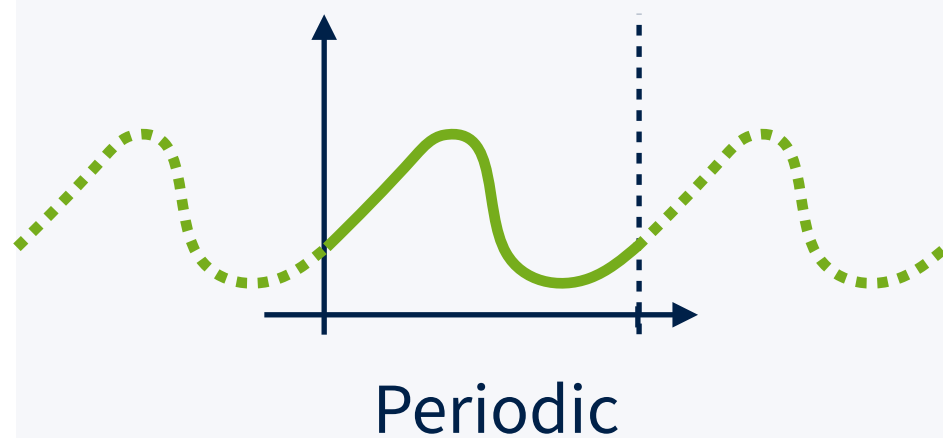
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Periodic or symmetric boundary condition

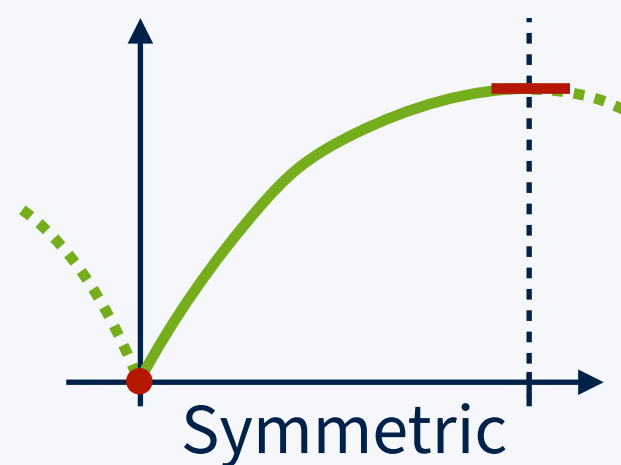
$$\nabla^2 \Psi = -f \Leftrightarrow \psi = G \star f \Leftrightarrow \hat{\psi} = \hat{G} \hat{f}$$

where $\hat{G} = -1/k^2$ is the Green's function in a spectral domain



$$f(x) \rightarrow \tilde{f}(k)$$

$\rightarrow$ 1D Discrete Fourier transform (real to complex or complex to complex)



$$f(x) \rightarrow \tilde{f}(k)$$

$\rightarrow$ 1D Discrete cosine/sine transform depending on the symmetry

... with various boundary conditions

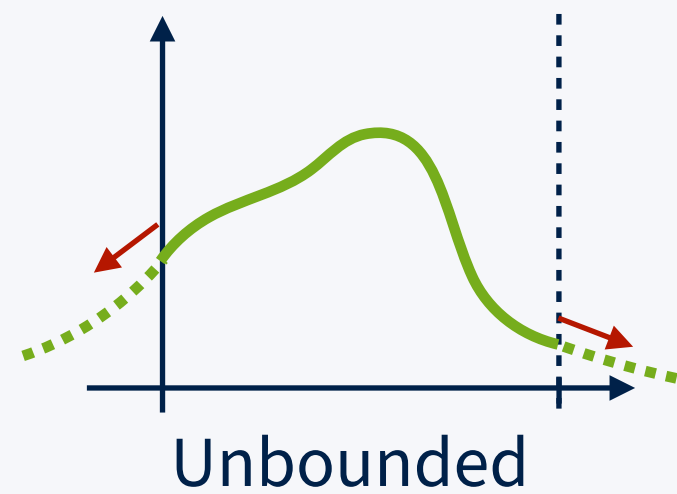
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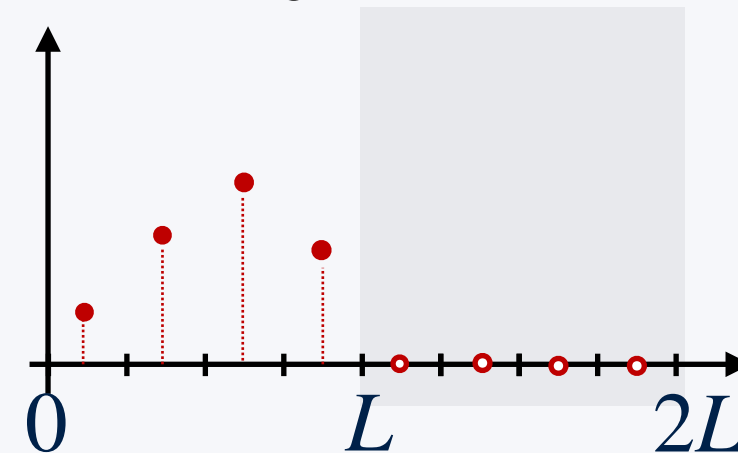
Unbounded and semi-unbounded boundary condition

$$\nabla^2 \Psi = -f \Leftrightarrow \psi = G_\delta \star f \Leftrightarrow \hat{\psi} = \hat{G}_\delta \hat{f} \quad \text{where } \hat{G}_\delta \text{ is the regularized Green's function}$$



$$f(x) \rightarrow \tilde{f}(k)$$

→ Domain doubling technique^[2,3]:



- f is extended to a domain of $2N$ with 0 padding
- Discrete Fourier transform on the padded function

... with various boundary conditions

- FFT-based method:

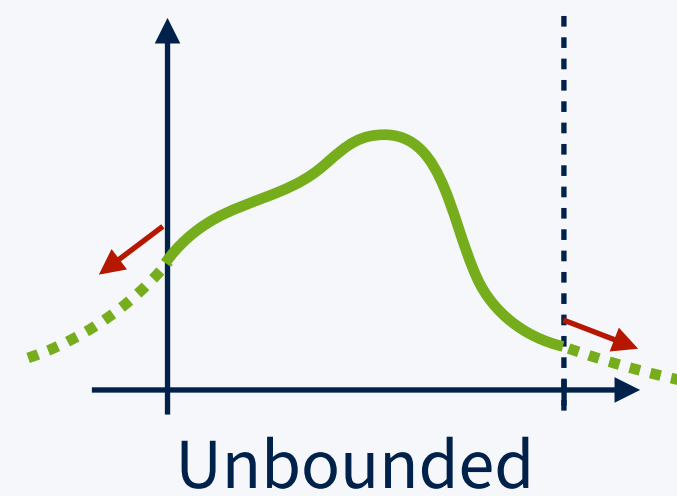
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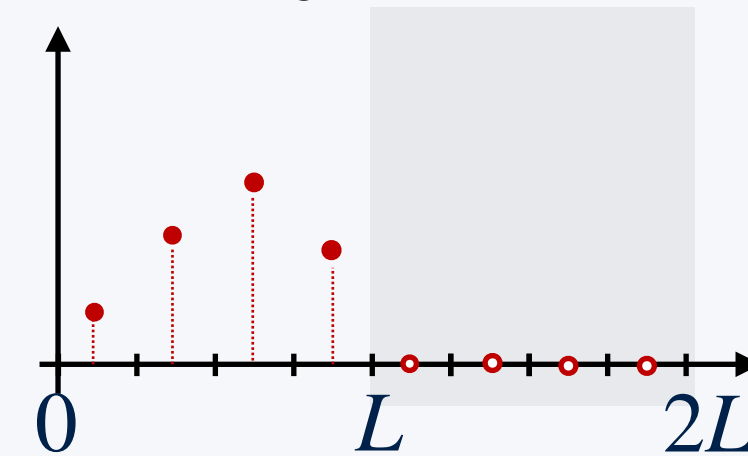
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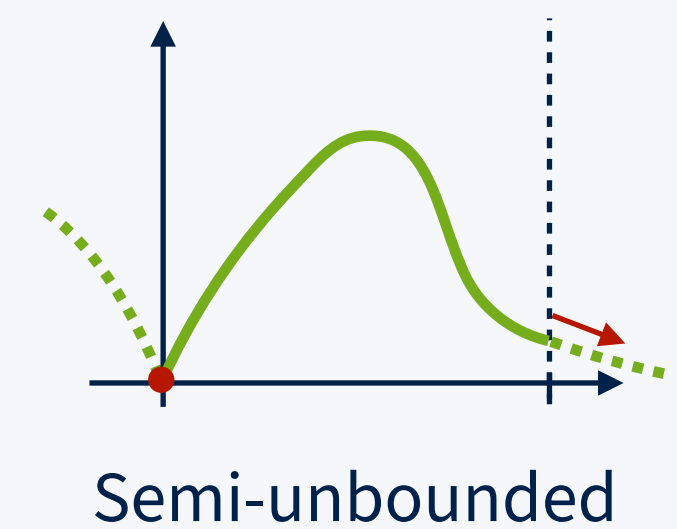


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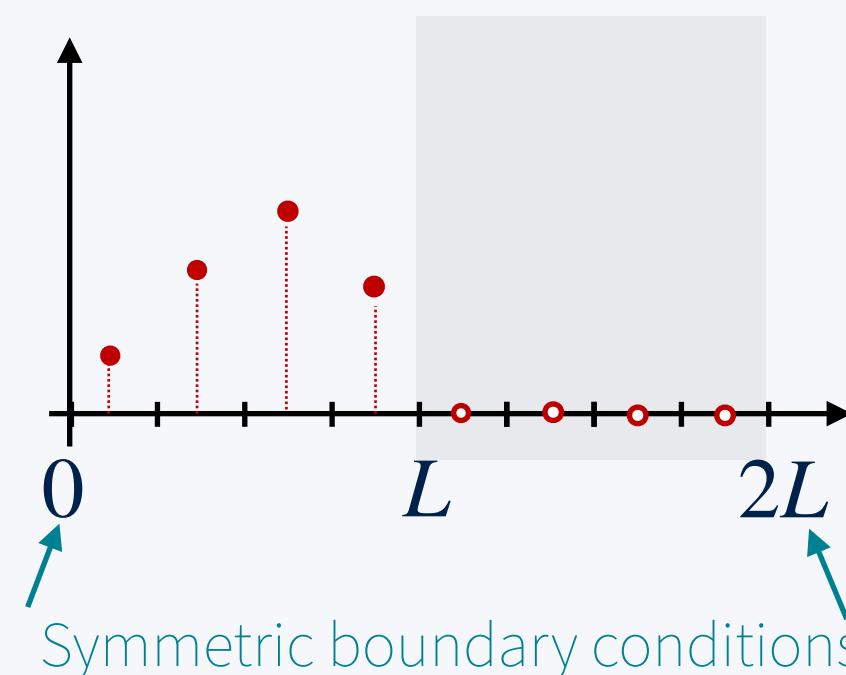
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$$f(x) \rightarrow \tilde{f}(k)$$



- f is extended to a domain of $2N$ with 0 padding
- Discrete sine/cosine transform on the padded function
- DST/DCT is equivalent to impose symmetric boundary conditions on the new domain

Very convenient for various problems:

- Inflow - outflow
- Wall bounded (wake in ground effects)

[2] Hockney:1988

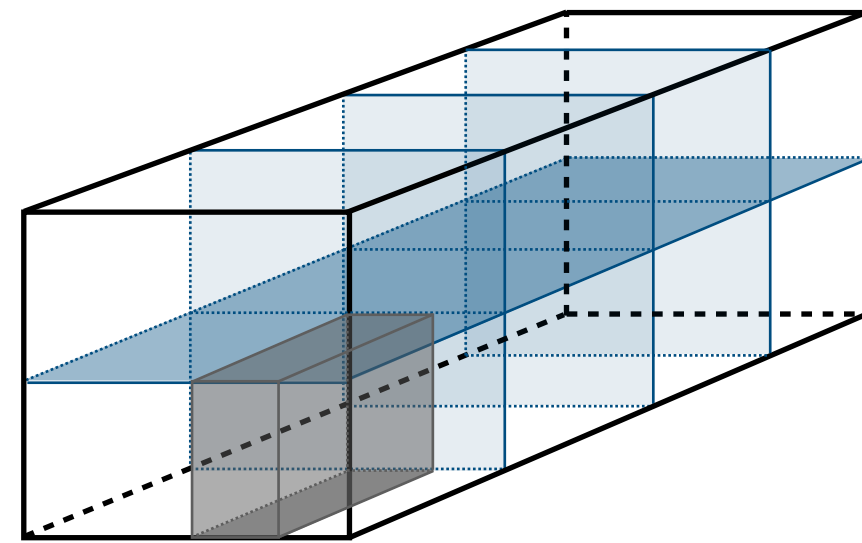
[3] Caprace : 2021

3D FFT on massively parallel systems

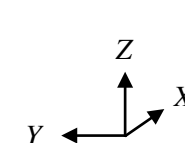
- Pencil decomposition
- The topology switches:

Source topology

	4		3
2			1



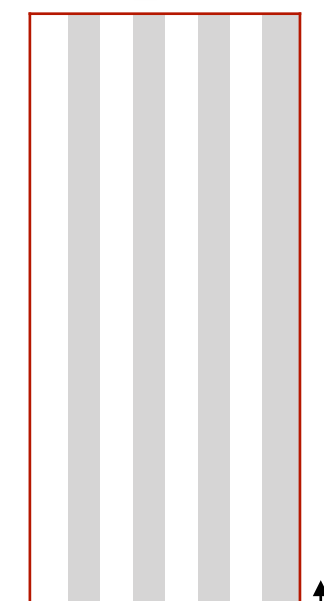
copy to buffer



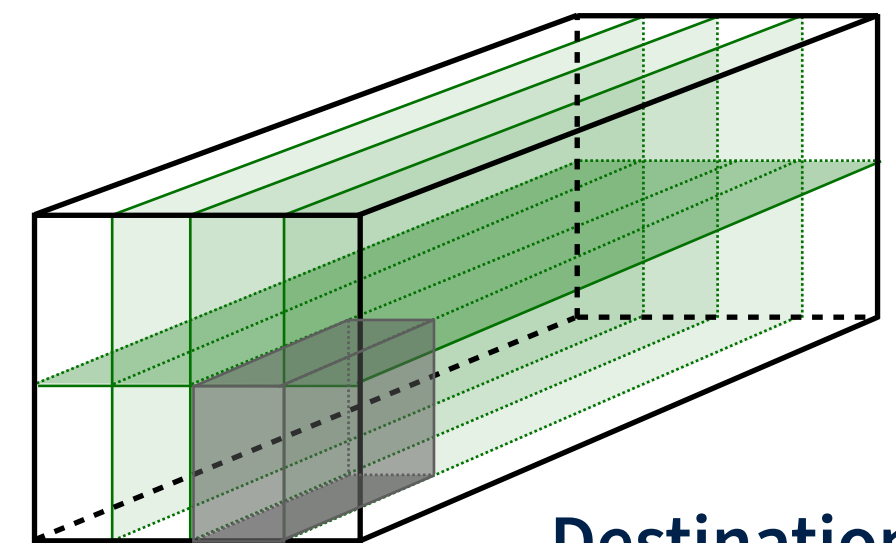
Communication



Shuffle



Copy back



Destination topology

8	7	6	5
4	3	2	1

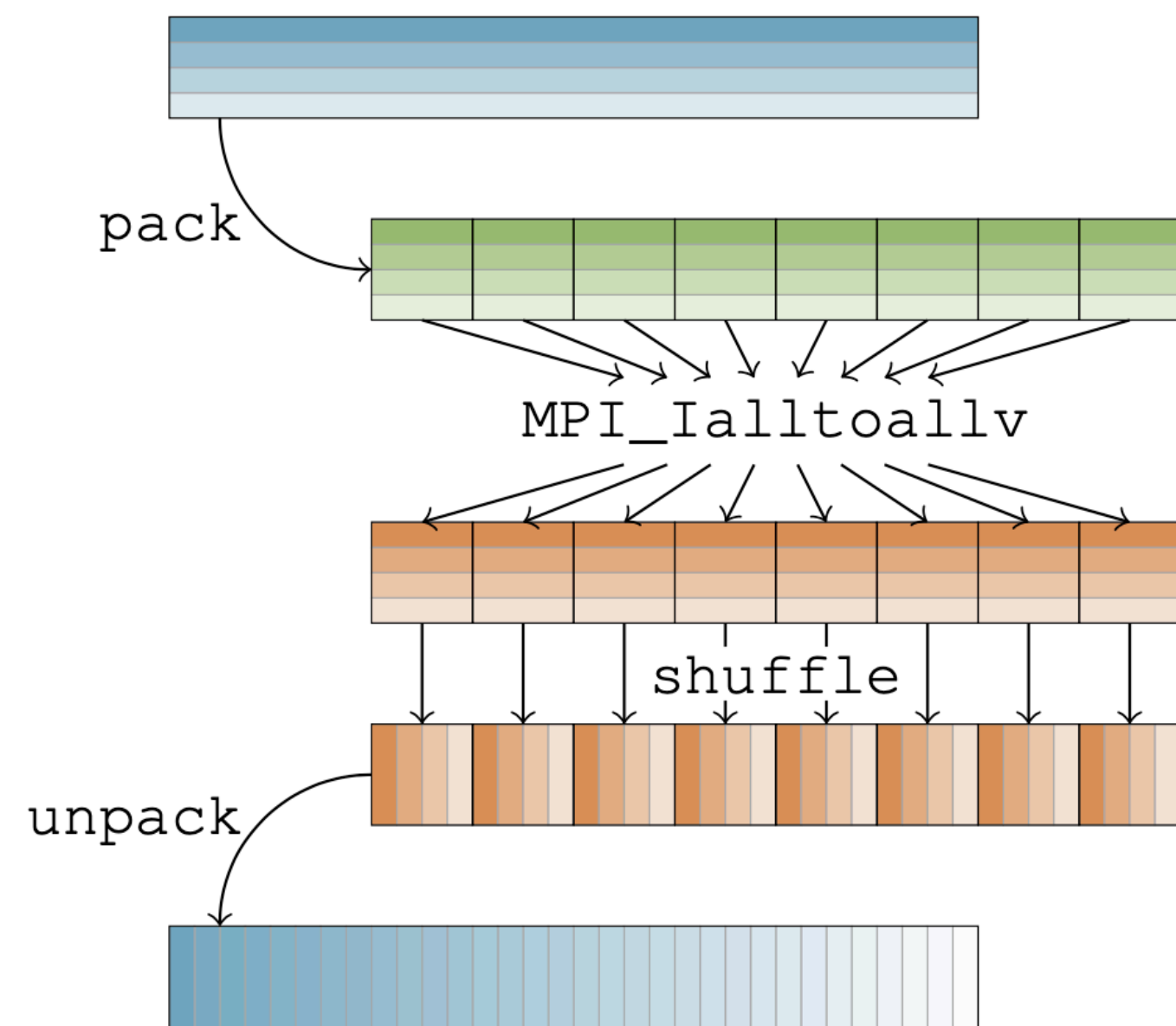
Distributed 1D FFTs require:

- All data on the same processor must be aligned in the direction of the FFT
- Unit-stride in memory

- An all-to-all communication problem - 3 implementations:
 - All-to-all communication
 - Non-blocking communication with manual packing and un packing
 - Non-blocking communication with MPI_Datatypes
- Optimizations:
 - Order of the 1D FFTs determined to reduce the memory footprint and the computational cost
 - Creation of sub-communicators to reduce the memory footprint on large partition

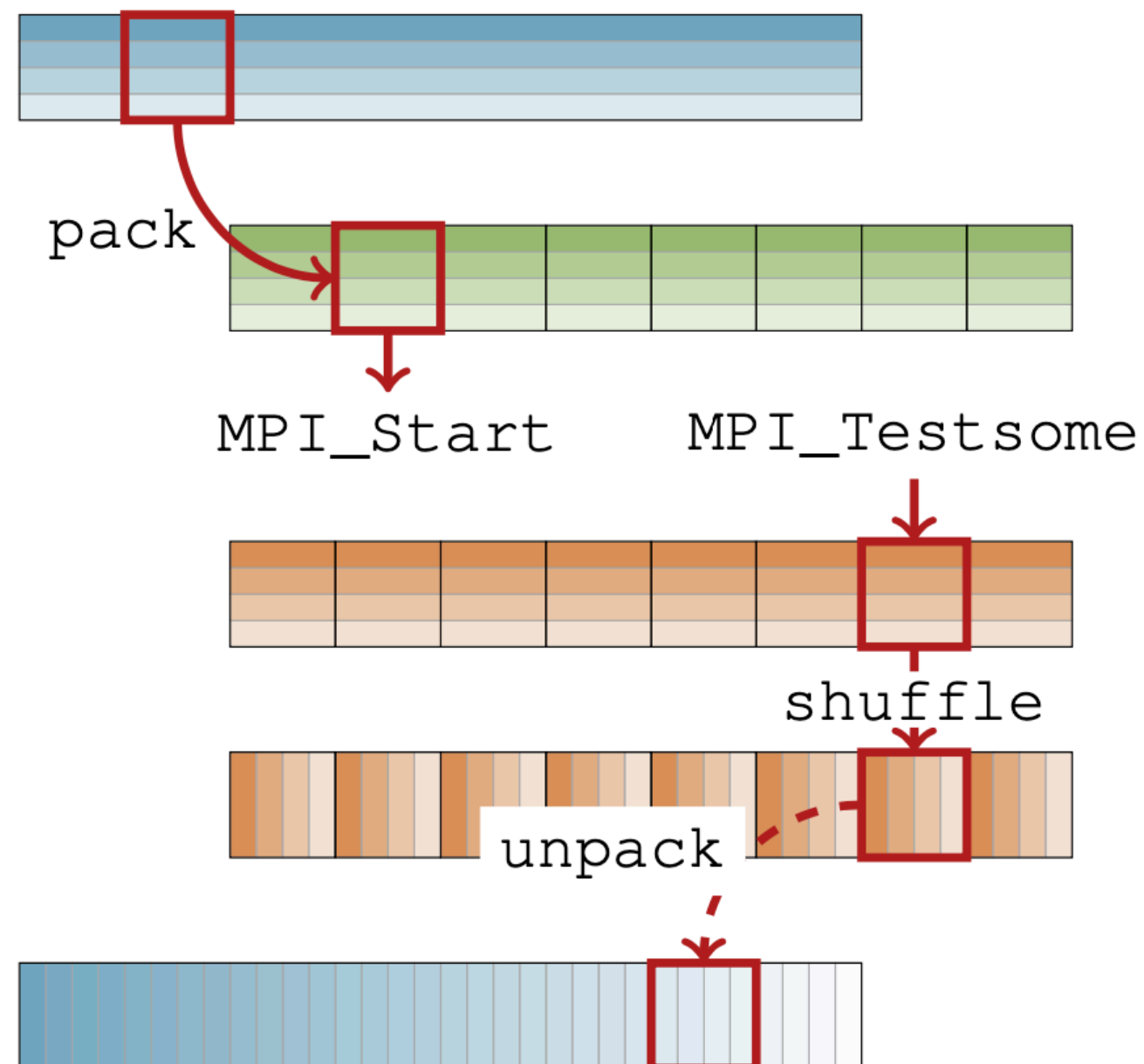
Switching between pencils - 3 implementations

All-to-all



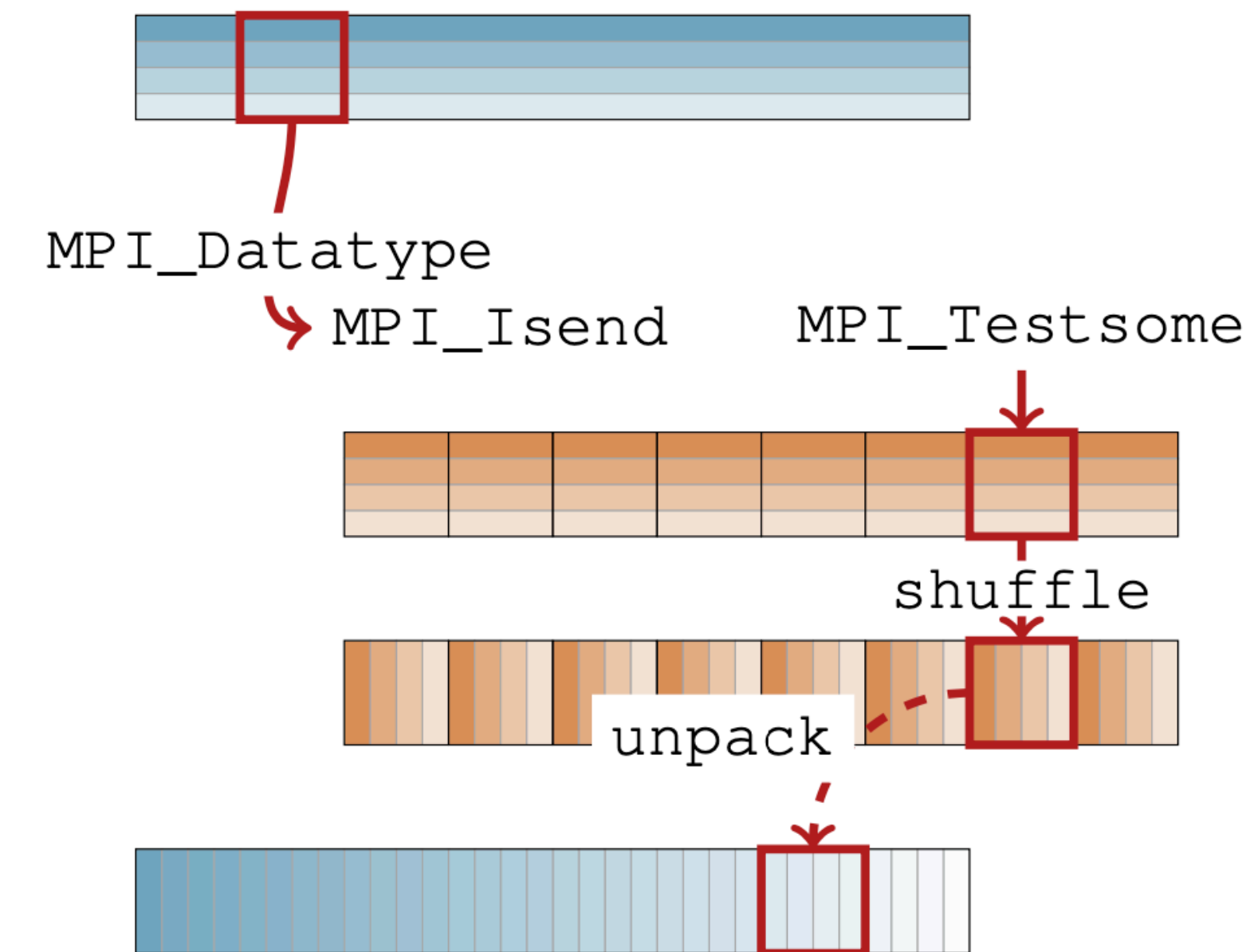
- **Simpler implementation** using MPI_Ialltoallv
- **Implicit barrier**
- Transposition of the data (based on FFTW)

Non-blocking with manual packing



- Use **persistent MPI_Send_Init, MPI_Recv_Init** and **MPI_Testsome**
- **Overlap** the **data packing** and the **transposition** with the **communication**
- Transposition of the data based on FFTW

Non-blocking with MPI_Datatype



- Use **non-blocking Send/Recv request**
- **Avoid manual packing** by using **MPI_Datatype**
- **Overlap** the **data packing** and the **transposition** with the **communication**
- Transposition of the data based on FFTW

1D FFTs reordering

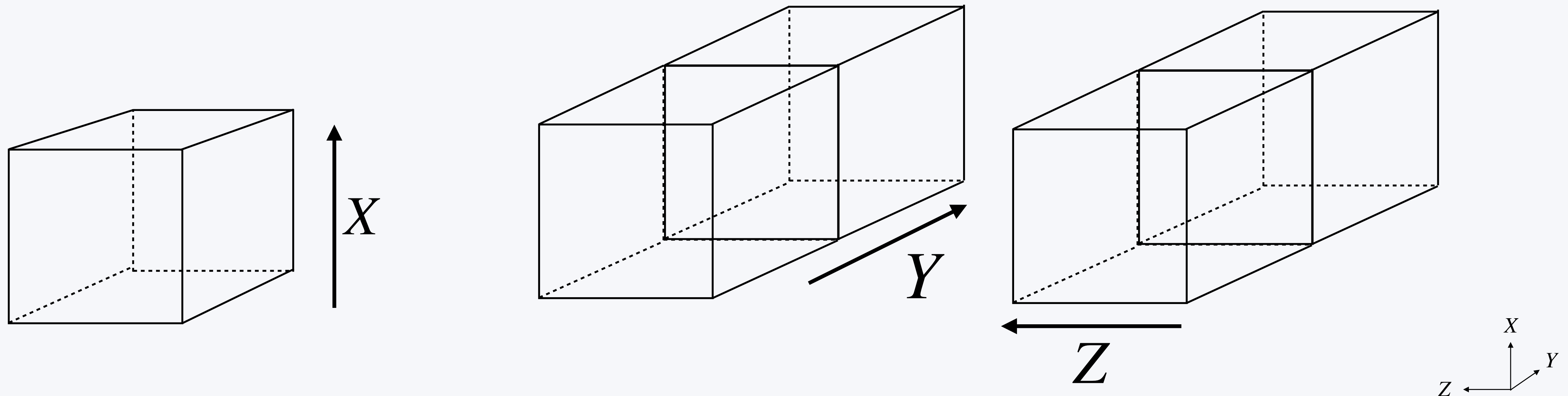
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Example 1 - (Periodic, Unbounded, Periodic)

Three 1D FFTs have to be performed:

Without reordering of the transform $X \rightarrow Y \rightarrow Z$:

- 1) $\psi_{x,y,z} \rightarrow \tilde{\psi}_{k_x,y,z}$ A real-to-complex DFT
- 2) $\tilde{\psi}_{k_x,y,z} \rightarrow \tilde{\psi}_{k_x,k_y,z}$ A complex-to-complex DFT on an extended and padded domain
- 3) $\tilde{\psi}_{k_x,k_y,z} \rightarrow \hat{\psi}_{k_x,k_y,k_z}$ A complex-to-complex DFT (on an extended domain)



→ FFT in Z costs $2N^3 \log(N)$

1D FFTs reordering

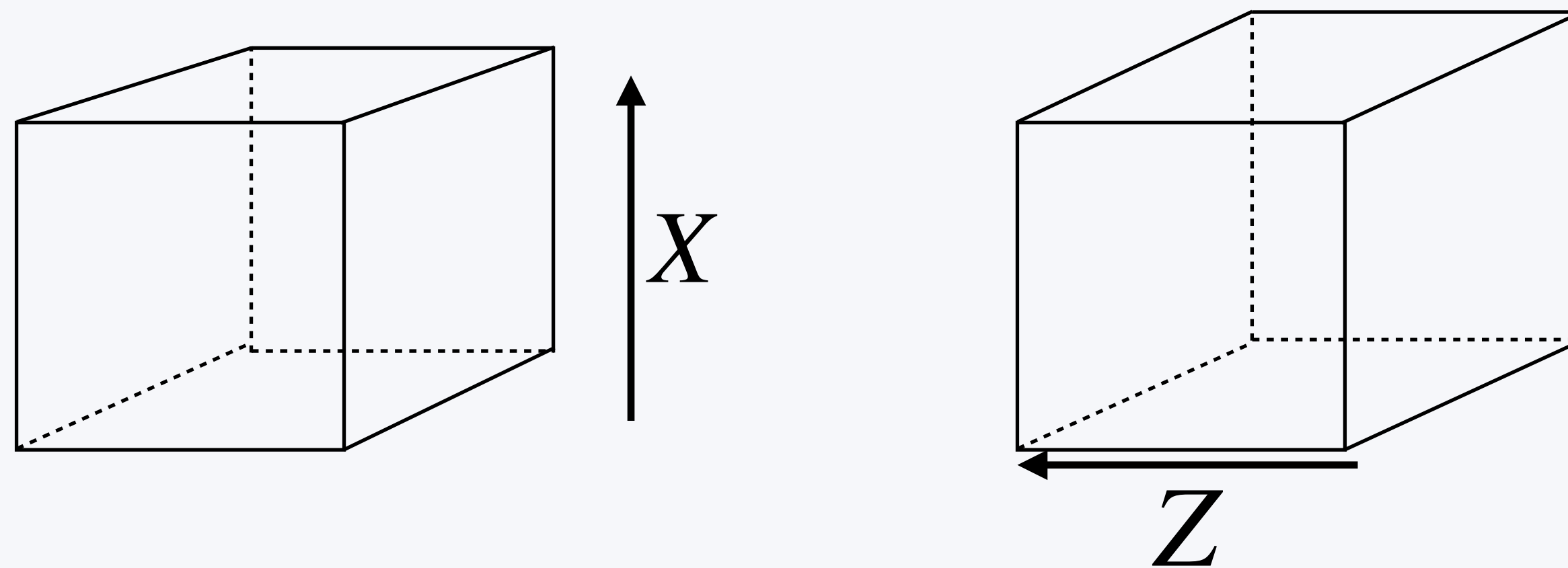
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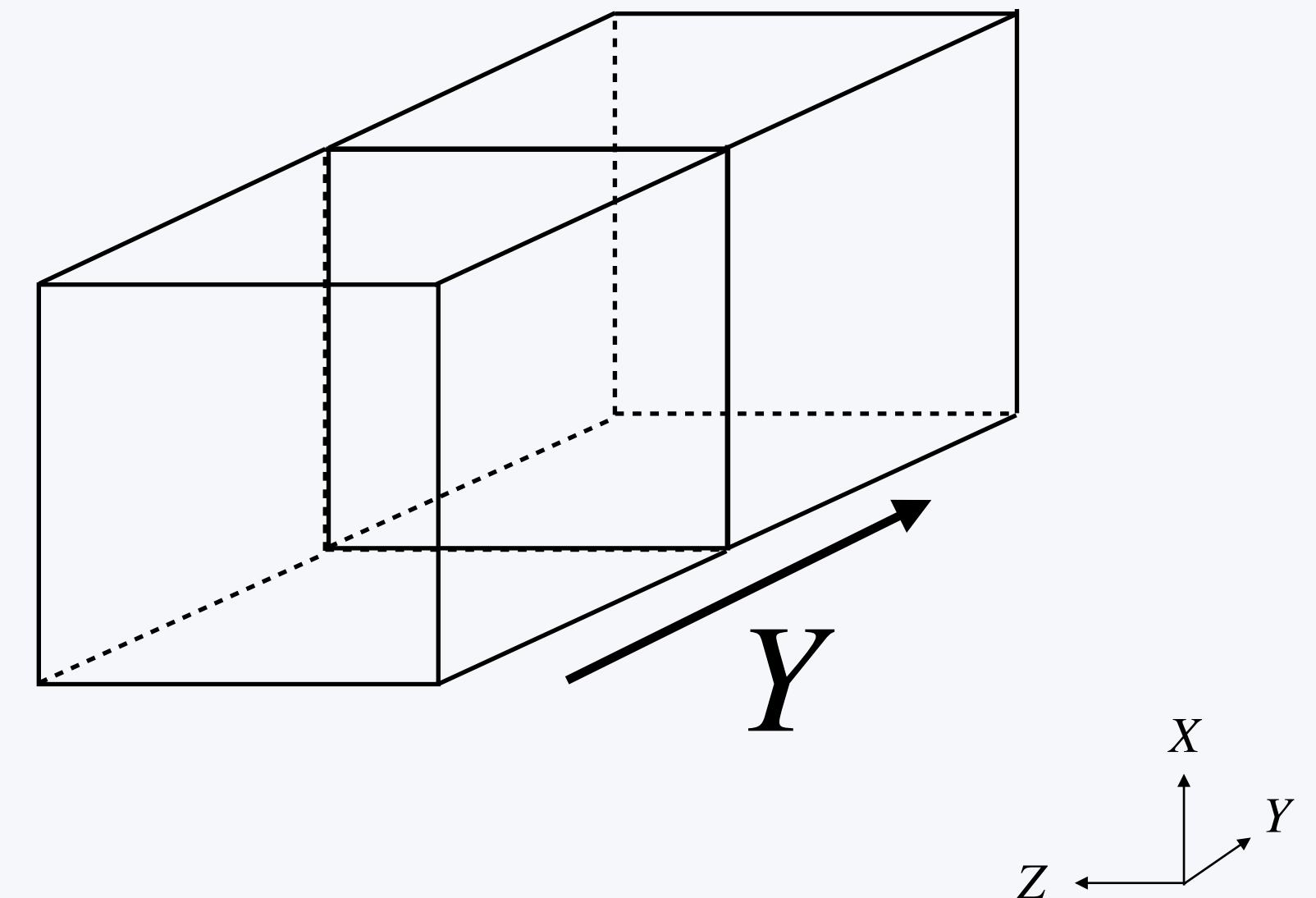
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- 3) $\tilde{\psi}_{k_x,y,k_z} \rightarrow \hat{\psi}_{k_x,k_y,k_z}$ A complex-to-complex DFT on an extended domain



→ FFT in Z costs $N^3 \log(N)$



Applications - Biot-Savart solver: Electromagnetism, Fluid mechanics,...

Methodology

- I. Forward transform of the rhs: $\omega \rightarrow \hat{\omega}$
- II. Computation of the curl in the spectral space
- III. Multiplication with the spectral Green's function $\hat{G}$

Testcase

- Compact vortex tube: $\omega(x, y, z) = \{0, 0, -\omega_z(r)\}$

$$\text{with } \omega_z(r) = \begin{cases} \frac{1}{2\pi} \frac{2}{R^2} \frac{1}{E_2(1)} \exp\left(-\frac{1}{\left(1 - \left(\frac{r^2}{R^2}\right)\right)}\right) & \text{if } r \leq R \\ 0 & \text{otherwise} \end{cases}$$

- Cubic domain of size $[0, L]^3$
- X and Y : fully unbounded
- Z : periodic

- Corresponding analytical velocity:

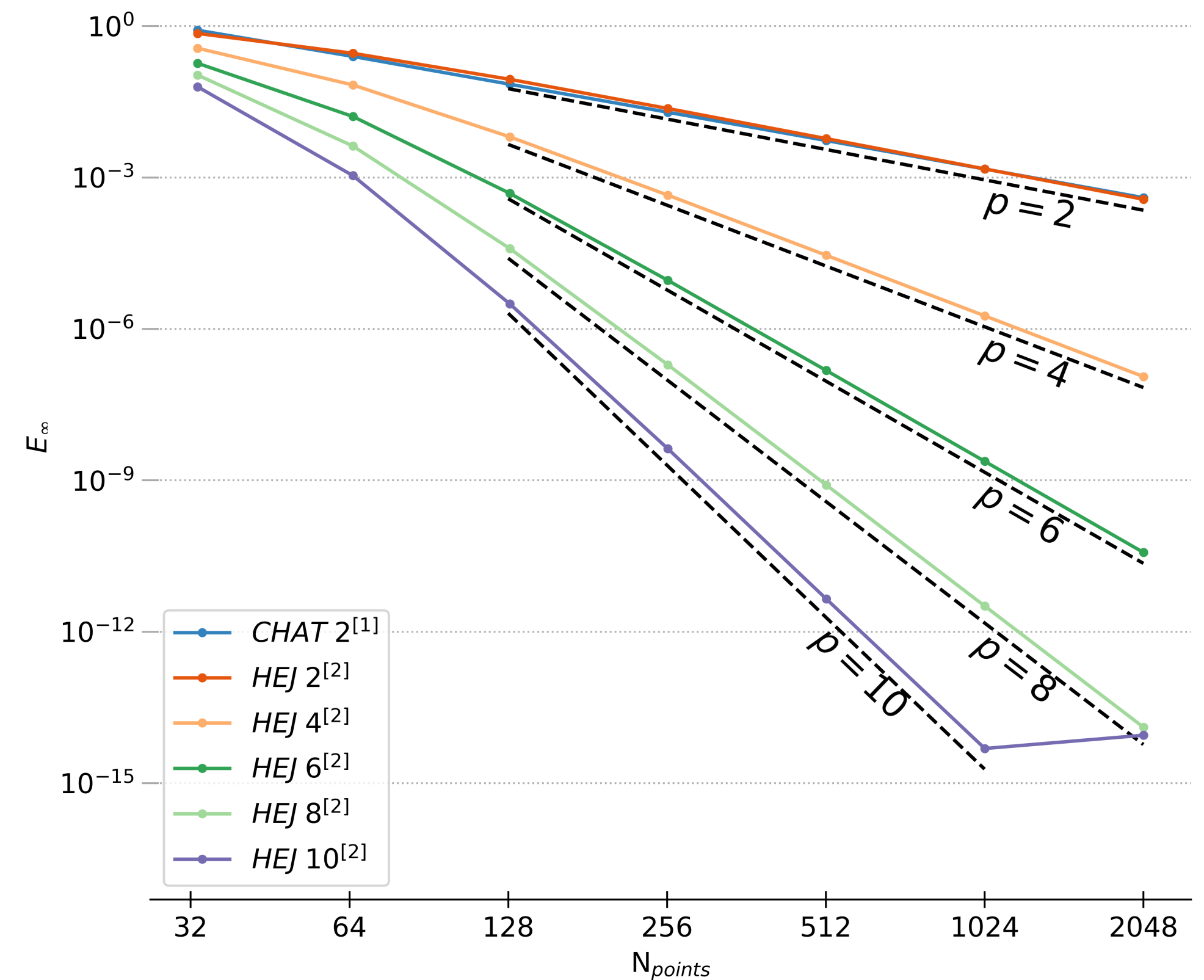
$$\mathbf{u}(x, y, z) = \left\{ -\sin(\theta)u_\theta(r), \cos(\theta)u_\theta(r), 0 \right\}$$

$$\text{with } u_\theta(r) = \begin{cases} \frac{1}{2\pi r} \left[1 - \frac{1}{E_2(1)} \left(1 - \left(\frac{r^2}{R^2} \right) \right) E_2\left(\frac{1}{1 - \left(\frac{r^2}{R^2} \right)} \right) \right] & \text{if } r \leq R \\ \frac{1}{2\pi r} & \text{otherwise} \end{cases}$$

$$\nabla^2 \mathbf{u} = -\nabla \times \omega$$

$$E_\infty = \sup_{x,y,z} \{ |\mathbf{u}(x, y, z) - \mathbf{u}_{ref}(x, y, z)| \}$$

Infinite norm of the error



[1] Green's function from Chatelain et al., 2010

[2] Green's function from Hejlesen et al., 2013

Results: Comparison with accFFT^[4]

Comparison with accFFT, one of the fastest distributed FFT libraries on CPU^[5]

Testcase

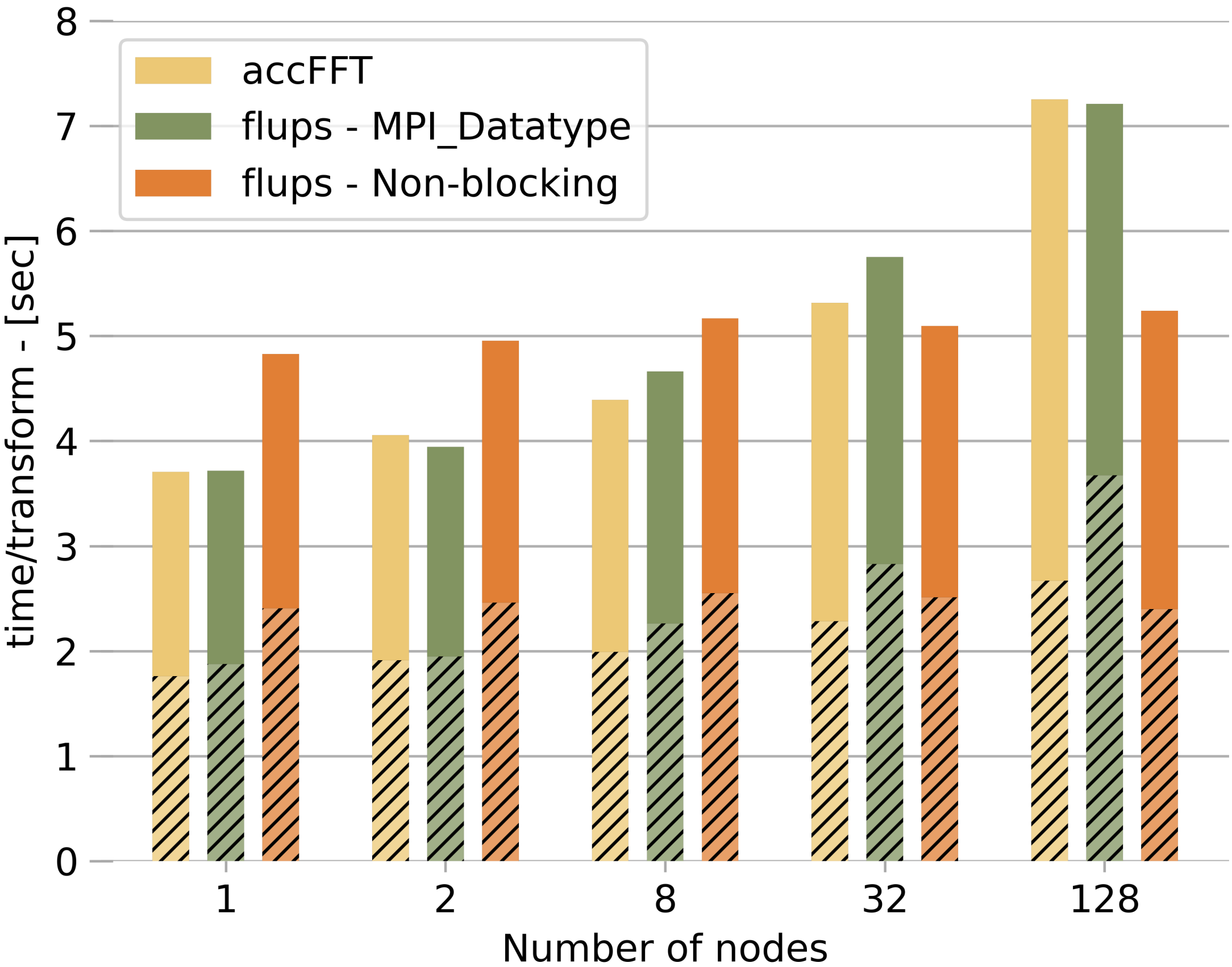
- Weak scaling
- Fully periodic
- Rectangular domain
- unknowns per rank: 256^3

accFFT

- Pencils are aligned in the Z direction
- Does $Z \rightarrow Y \rightarrow X$
- Opt. Flag: ACCFFT_MEASURE

Flups

- Pencils are aligned in the X direction
- Does $X \rightarrow Y \rightarrow Z$
- Opt. Flag: FFTW_MEASURE



MeluXina:

- CPU: AMD EPYC 7H12
- Interconnect: 200 Gb/s Infiniband HDR
- MPI: MPICH 4.1a1
- Transport Layer: UCX 1.13.1

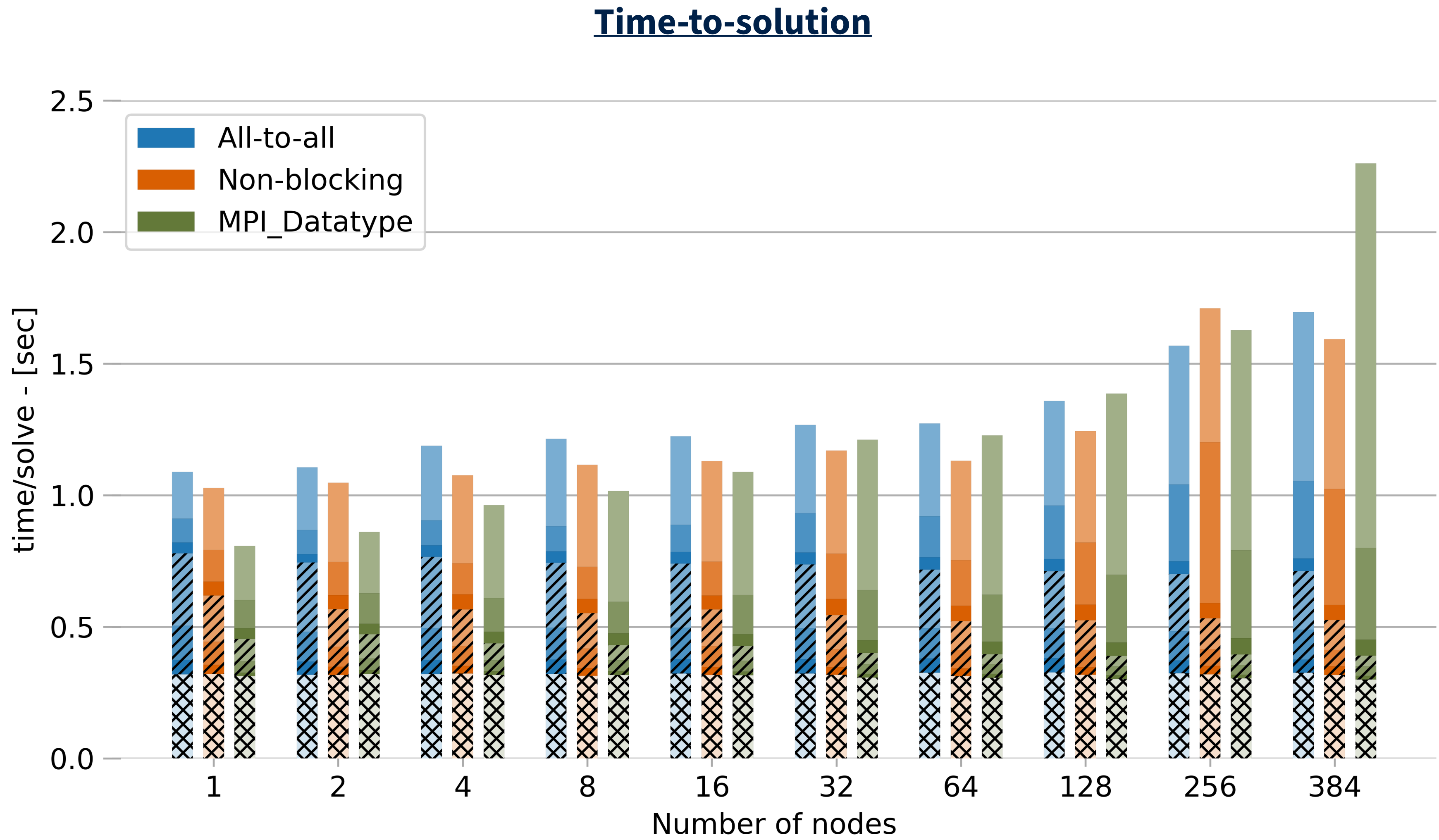
	Px	Py	Pz
1	1	8	16
2	1	16	16
8	1	32	32
32	1	64	64
128	1	128	128

→ At 128 nodes, non-blocking version of flups is 27% faster than accFFT

[4] Gholami:2015

[5] Ayala:2021a

Weak scaling - from 128 to 49,152 processes



Process distribution

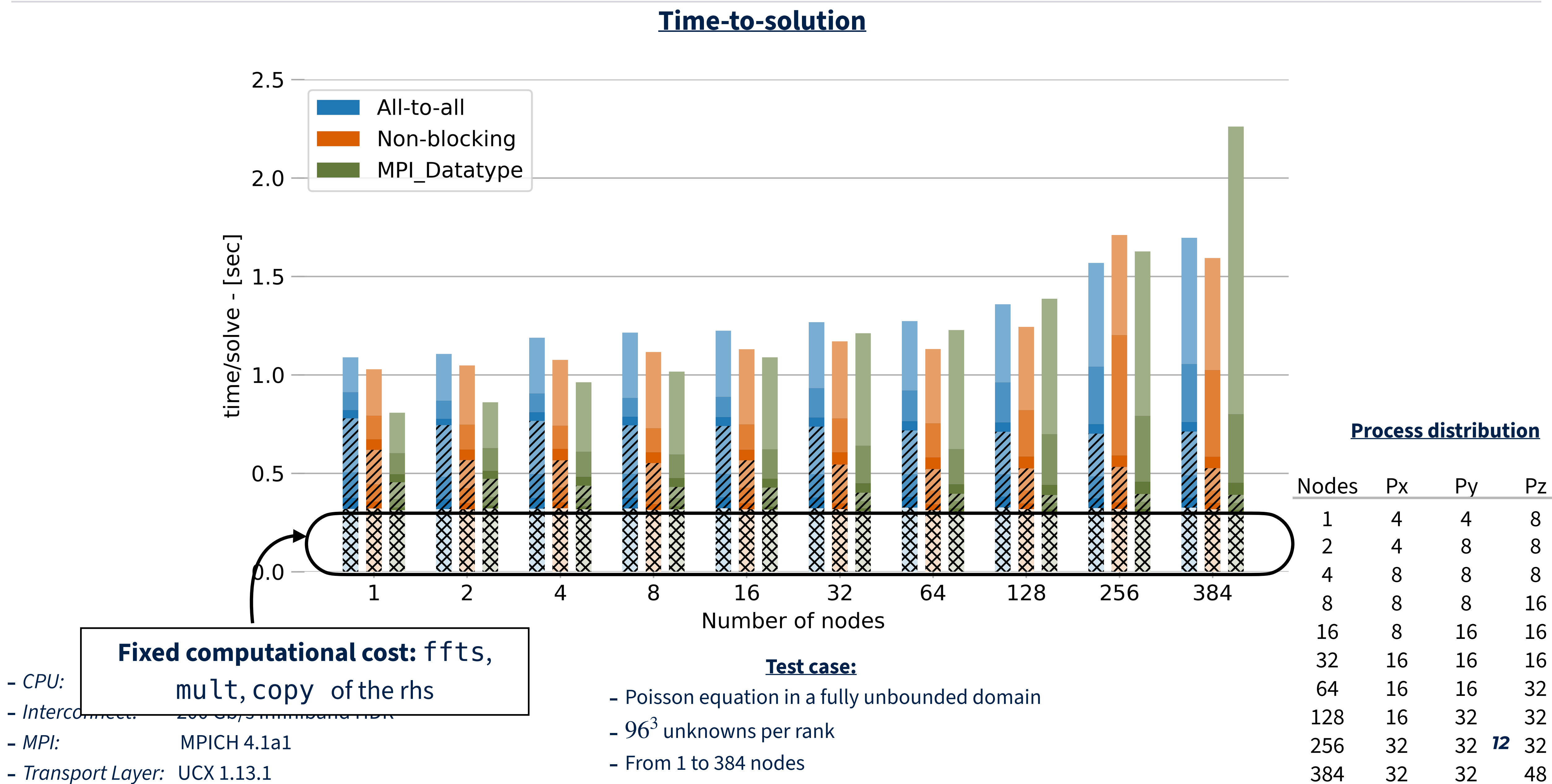
Nodes	Px	Py	Pz
1	4	4	8
2	4	8	8
4	8	8	8
8	8	8	16
16	8	16	16
32	16	16	16
64	16	16	32
128	16	32	32
256	32	32	32
384	32	32	48

- MeluXina:**

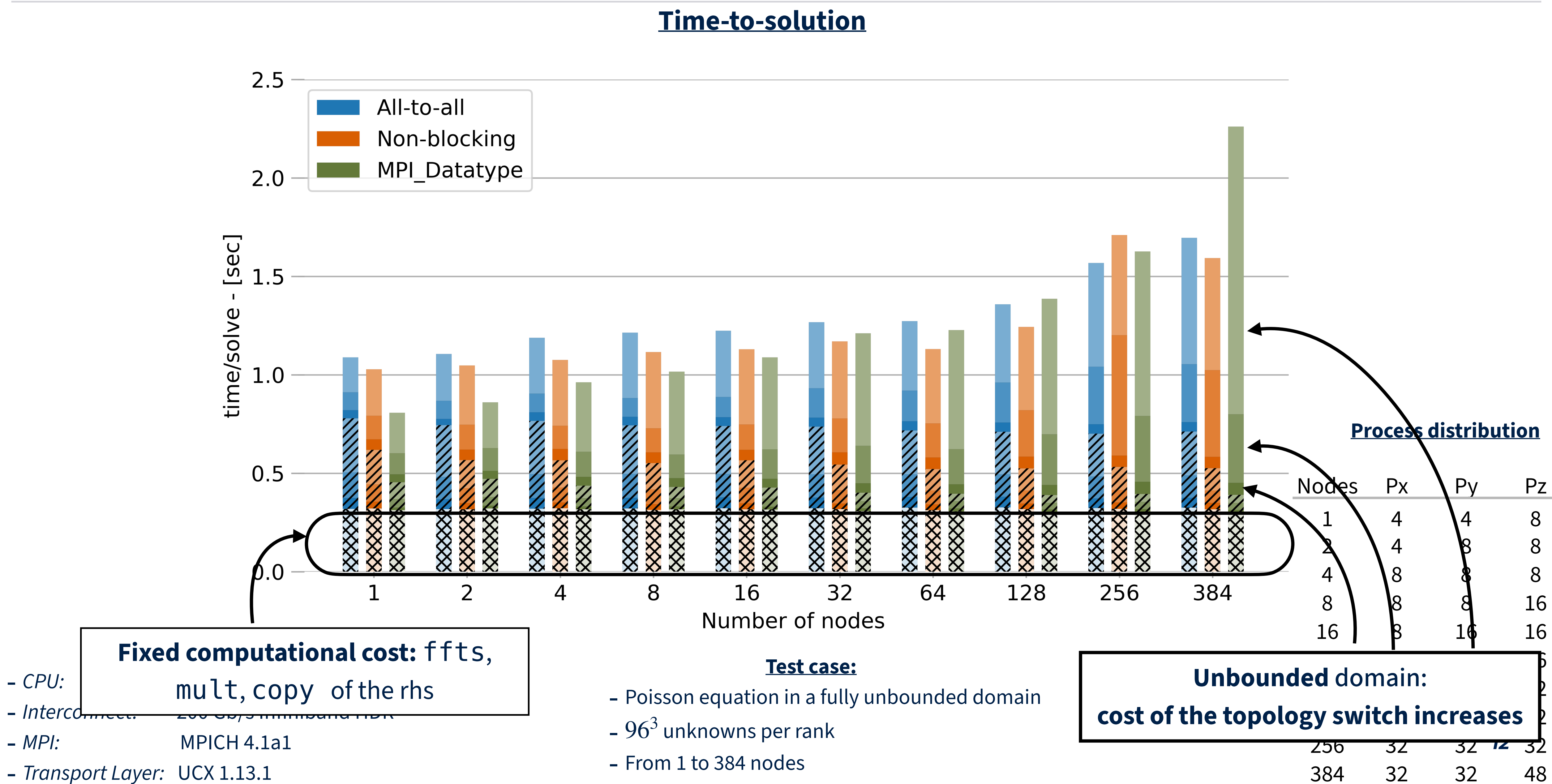
 - CPU: AMD EPYC 7H12
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- Test case:**

 - Poisson equation in a fully unbounded domain
 - 96^3 unknowns per rank
 - From 1 to 384 nodes

Weak scaling - from 128 to 49,152 processes

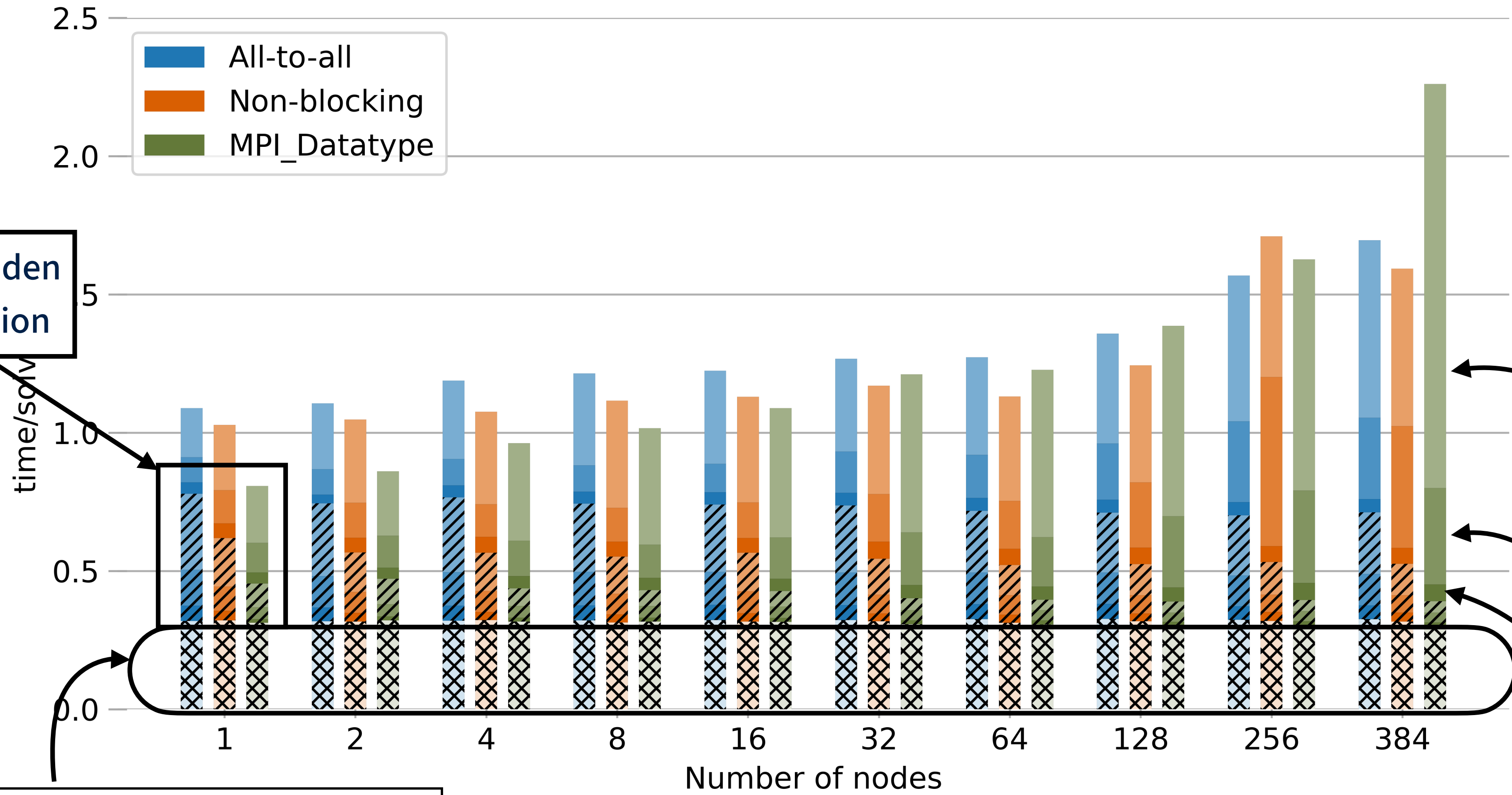


Weak scaling - from 128 to 49,152 processes



Weak scaling - from 128 to 49,152 processes

Time-to-solution



Process distribution

Nodes	Px	Py	Pz
1	4	4	8
2	4	8	8
4	8	8	8
8	8	8	16
16	8	16	16
256	32	32	12
384	32	32	48

Manual packing hidden in the communication

Fixed computational cost: ffts, mult, copy of the rhs

- CPU: 200 GB/s memory bandwidth
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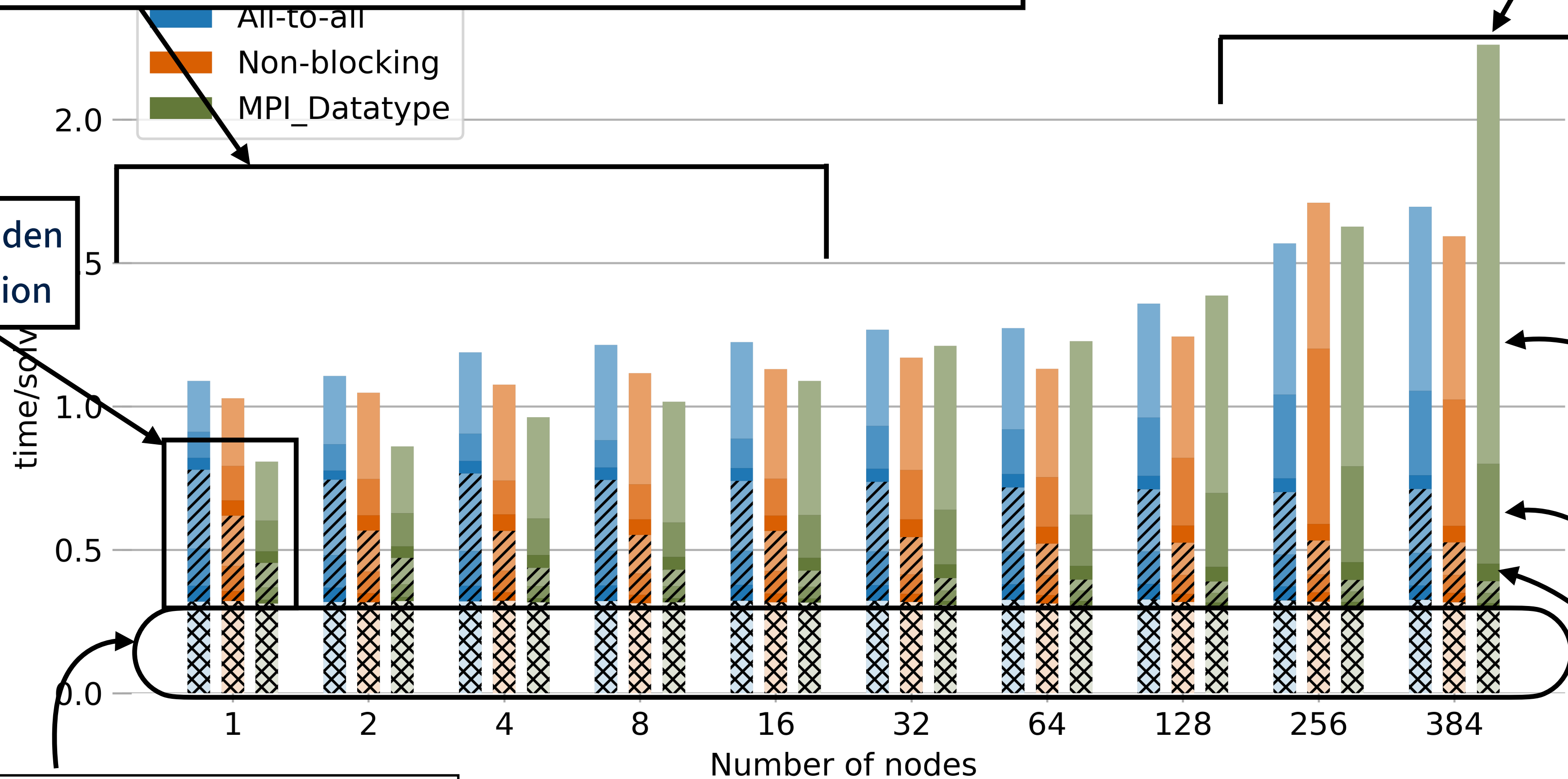
Unbounded domain: cost of the topology switch increases

Weak scaling - from 128 to 49,152 processes

Time-to-solution:
MPI_Datatype > all to all > Non-Blocking

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Manual packing hidden
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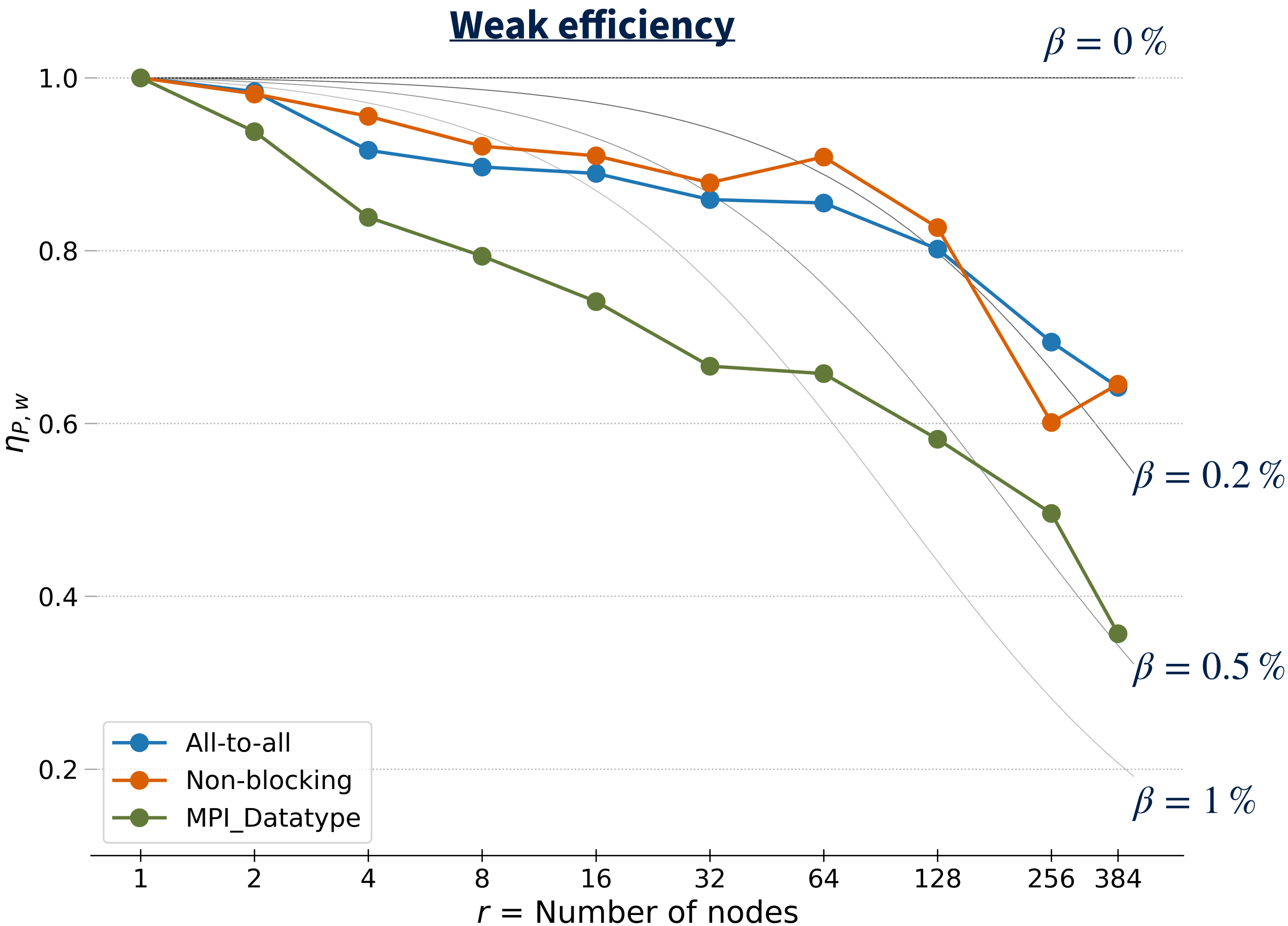
Gustafson's law:

efficiency η is

$$\eta_{P,w} = \frac{1}{1 + (r - 1)\beta}$$

where:

- $r = N/N_0$ is the ratio of the computational resource
- β is the sequential percentage of the program



Weak scaling - from 128 to 49,152 processes

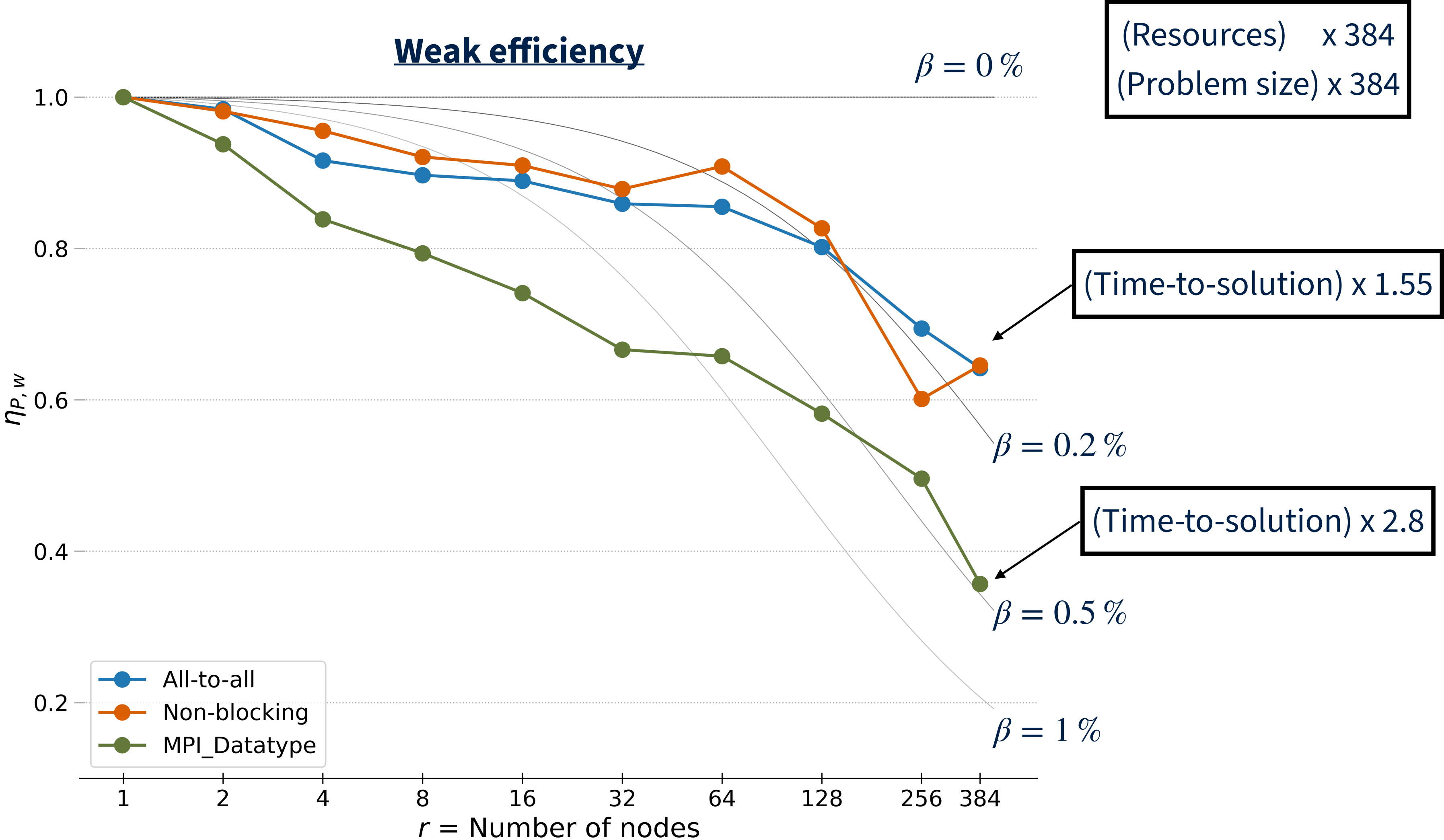
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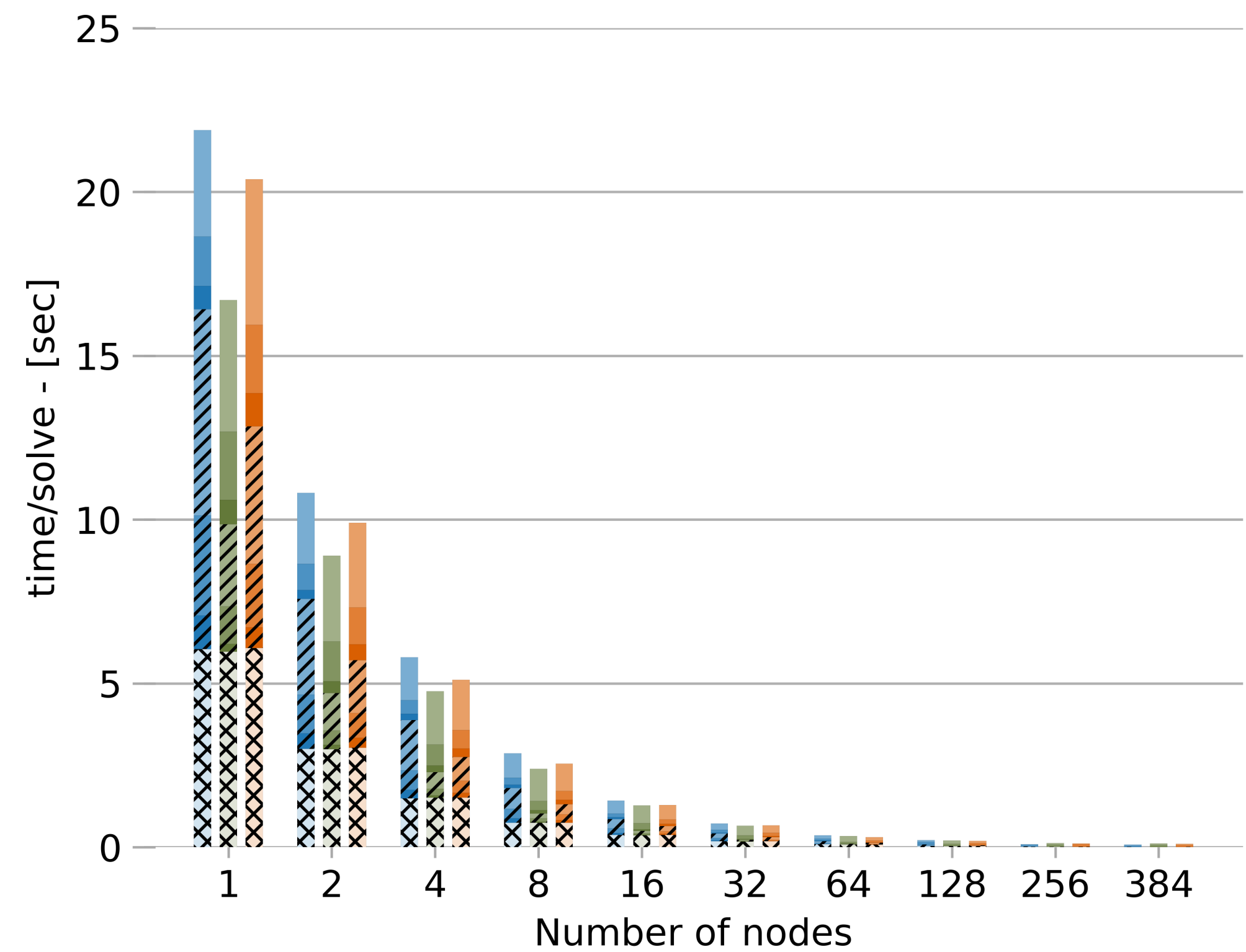
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Strong scaling - from 128 to 49,152 processes

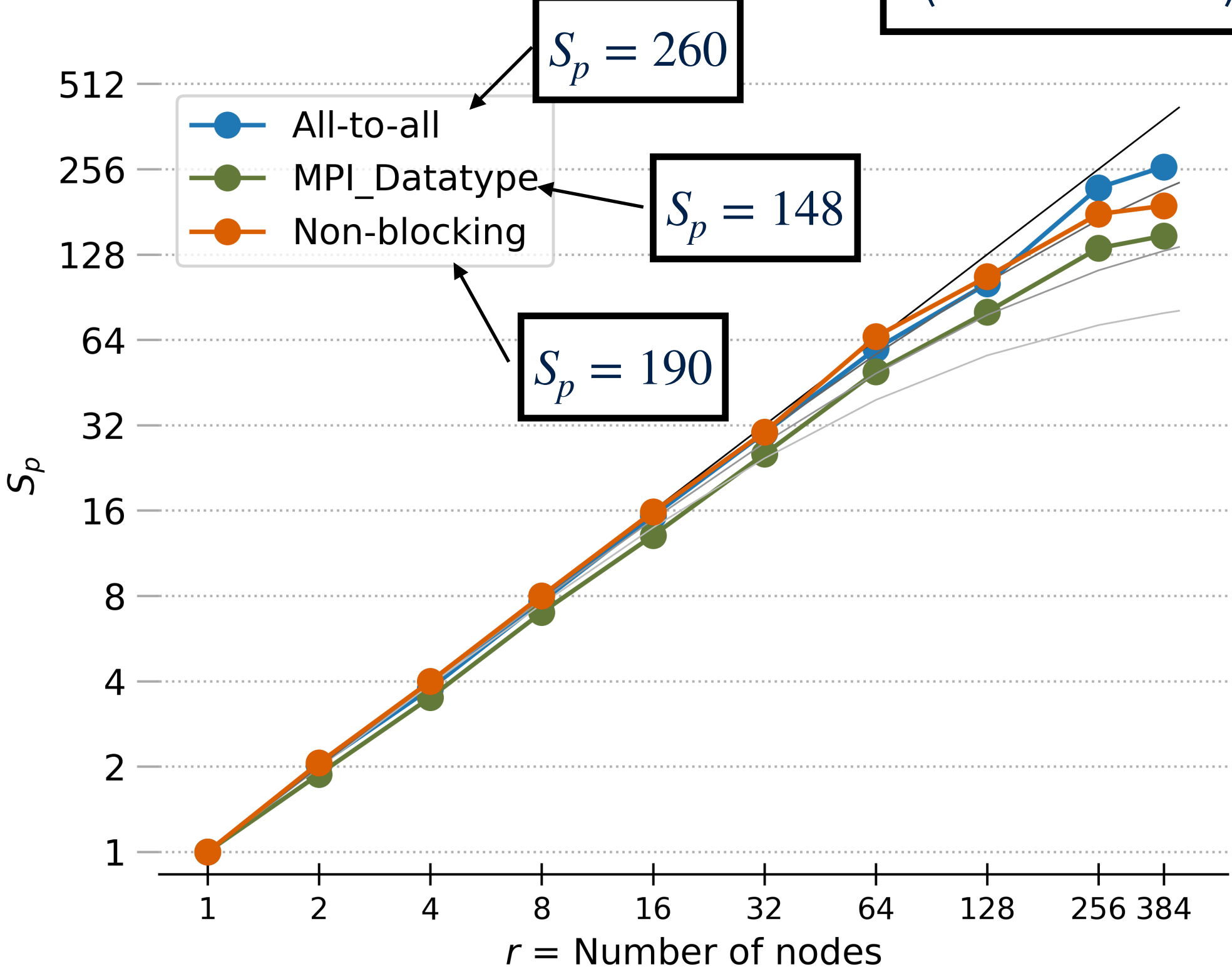
Time-to-solution



MeluXina:

- CPU: AMD EPYC 7H12
- Interconnect: 200 Gb/s Infiniband HDR
- MPI: MPICH 4.1a1
- Transport Layer: UCX 1.13.1

Speed up



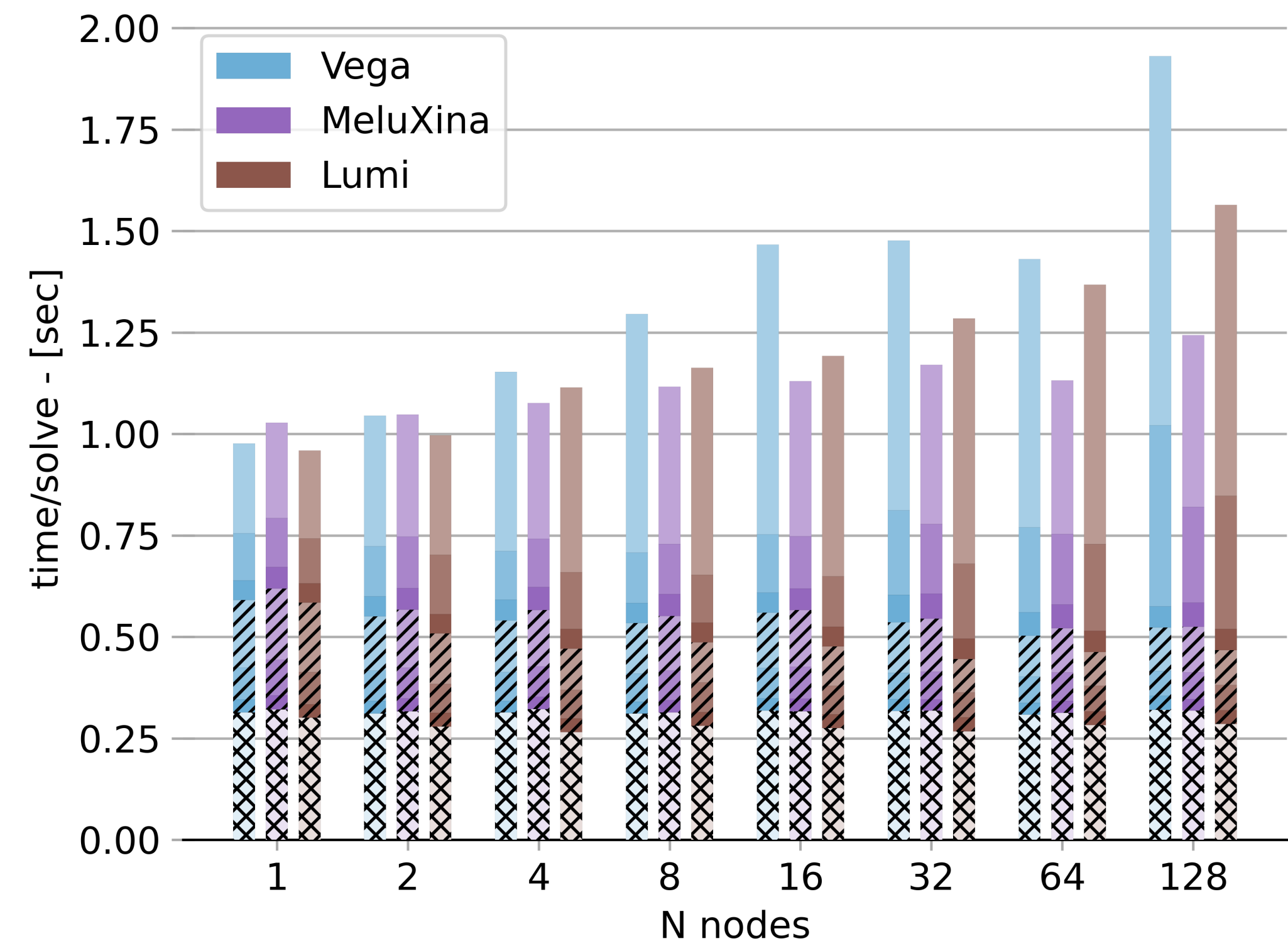
(Resources) x 384
(Problem size) x 1

Test case:

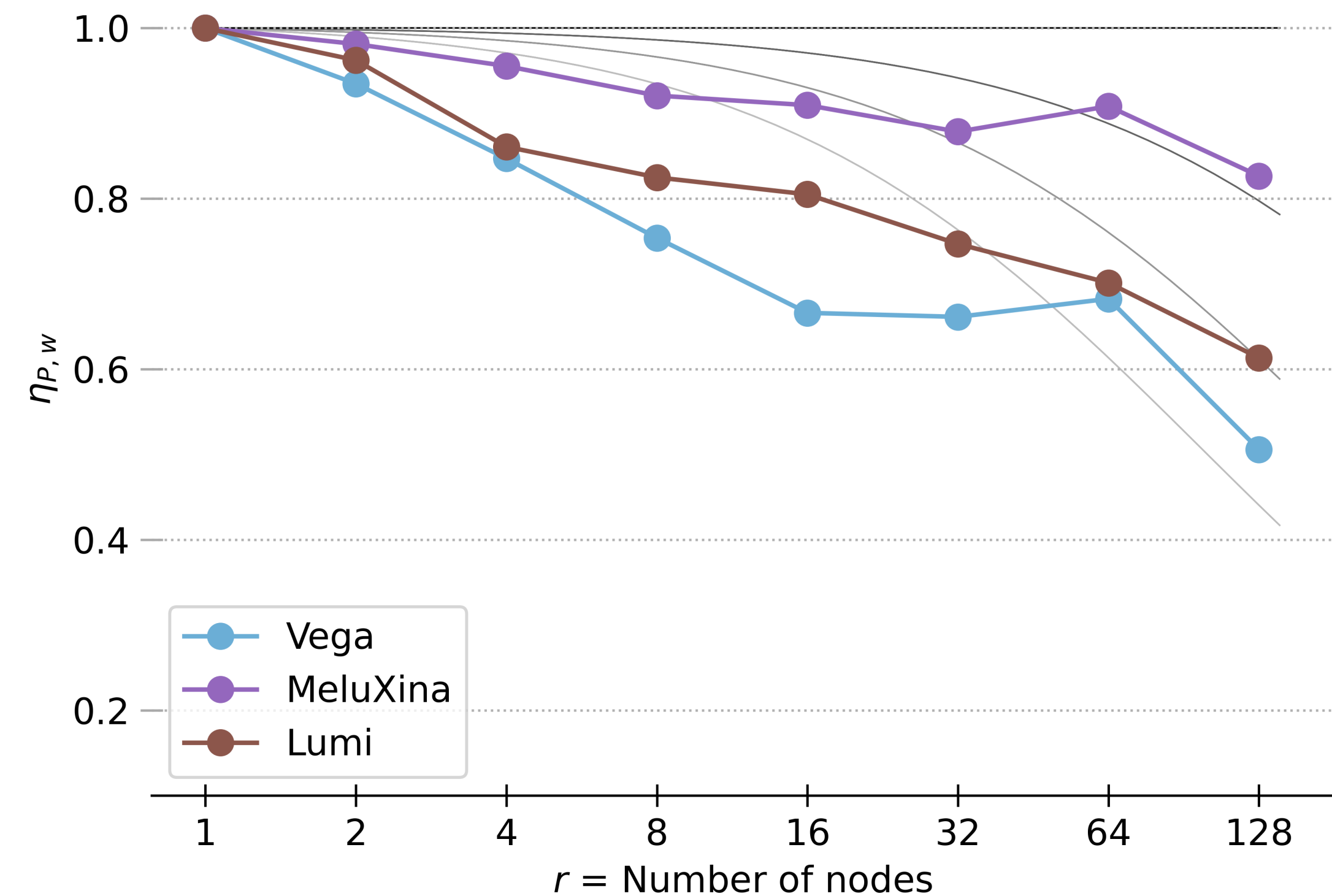
- Poisson equation in a fully unbounded domain
- 1281^3 unknowns
- From 1 to 384 nodes

European cluster comparison - Weak scaling test

Time-to-solution



Weak efficiency



Test case:

- Poisson equation in a fully unbounded domain
- 96^3 unknowns per rank
- From 1 to 128 nodes

Name	Location	CPU	Interconnect	Transport Layer	OSU Latency
Lumi	Finland	AMD EPYC 7763	200 Gb/s Slingshot-11	Libfabric 15.0.0	2.05 us
MeluXina	Luxembourg	AMD EPYC 7H12	200 Gb/s Infiniband HDR	ucx 1.13.1	1.45 us
Vega	Slovenia	AMD EPYC 7H12	100 Gb/s Infiniband HDR	ucx 1.13.1	1.99 us

Conclusions

Flups offers highly efficient distributed FFT-based Poisson solvers

Methodology

- Handle arbitrary cartesian topology
- 1000 combinations of boundary conditions
- 8 different Green's function
- 2 data layouts

Parallel performance

- Faster time-to-solution compared to accFFT on large FFTs
- Weak and Strong scalability with a parallel percentage $\approx 96\% - 98\%$
- Scalability on three EuroCC systems

Thanks for your attention!

Any questions?

- *FLUPS* - a flexible and performant massively parallel Fourier transform library, Balty et al., IEEE - Transactions on Parallel and Distributed Systems, 2023 (*forthcoming*)
- *FLUPS - A Fourier-based Library of Unbounded Poisson Solvers*, Caprace et al., SIAM Journal on Scientific Computing, vol. 43, no. 1, pp. C31–C60, January 2021
- <https://github.com/vortexlab-uclouvain/flups>

